Mixed-integer Programming based Differential and Linear Cryptanalysis

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Outline

□ Introduction

Mixed-Integer Programming (MIP) based Automatic Differential (Linear) Characteristic Search

Applications, Improvements, and Adaptation for Specific Ciphers

Future work: Domain Specific Language (DSL) for Constraint Programming based Automatic Cryptanalysis

Introduction

Differential cryptanalysis [1] is one of the most powerful attacks on block ciphers

- ✓ Truncated differential attack
- \checkmark Related-key differential attack
- ✓ Boomerang attack
- ✓ ...

Finding a good (related-key) differential (characteristic) with high probability is the first step in the (related-key) differential attack

[1] Eli Biham and Adi Shamir. Differential cryptanalysis of DES-like cryptosystems. Journal of Cryptology, 4(1):3-72, 1991.

Existing methods for characteristic search

Matsui's algorithm and its variants

- Mitsuru Matsui, On correlation between the order of S-boxes and the strength of DES, Eurocrypt 1994.
- Alex Biryukov, Ivica Nikolic.: Search for related-key differential characteristics in DES-like ciphers. FSE 2011
- \checkmark ... other branch and bound based algorithms

SMT(Satisfiability Modulo Theory) based Methods

- Nicky Mouha and Bart Preneel. Towards finding optimal differential characteristics for ARX: Application to Salsa20. IACR Cryptology ePrint Archive, Report 2013/328, 2013.
- ✓ Jean-Philippe Aumasson, Philipp Jovanovic, Samuel Neves. Analysis of NORX: Investigating Differential and Rotational Properties. In Latincrypt 2014, 2014.
- Stefan Kölbl and Gregor Leander and Tyge Tiessen. Observations on the SIMON block cipher family. CRYPTO 2015.

Integer programming based method

Existing methods for characteristic search

Integer programming based methods

- Nicky Mouha. Differential and linear cryptanalysis using mixed-integer linear programming. Information Security and Cryptology. Springer Berlin Heidelberg, 2012.
- ✓ Shengbao Wu, Mingsheng Wang. Security Evaluation against Differential Cryptanalysis for Block Cipher IACR ePrint 2011/551.
- Siwei Sun, Lei Hu, Peng Wang, Kexin Qiao, Xiaoshuang Ma, Ling Song. Automatic Security Evaluation and (Related-key) Differential Characteristic Search : Application to SIMON, PRESENT, LBlock, DES(L) and Other Bitoriented Block Ciphers. Asiacrypt 2014.

 Siwei Sun, Lei Hu, Meiqin Wang, Peng Wang, Kexin Qiao, Xiaoshuang Ma, Danping Shi, Ling Song: Towards Finding the Best haracteristics of Some Bit-oriented Block Ciphers and Automatic Enumeration of (Relatedkey) Differential and Linear Characteristics with Predefined Properties. IACR Cryptology ePrint Archive 2014:747 (2014) □ Mixed-integer programming (MIP), an example

- ✓ Objective function
- ✓ Constraints

 $\begin{array}{rl} \min -x_1 + x_2 - 2x_3 + x_4 - x_5 \\ \text{subject to} & & \\ x_1 + x_2 & \leq & 1 \\ x_1 - 5x_2 + x_3 & \leq & 2 \\ 2x_3 + 2x_4 - 4x_5 & \leq & 1 \\ x_2 - 2x_4 + x_5 & \leq & 0 \\ & x & \in & \{0, 1\}^5 \end{array}$

Feasible region: the set of all solutions satisfying the constraints

Mixed-integer programming based method

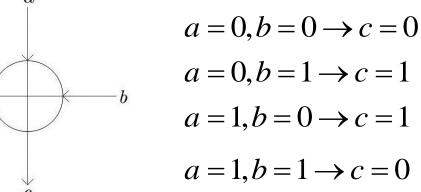
Basic Idea

- ✓ Describe the propagation characteristic of the difference patterns using linear inequalities
- ✓ Find solutions of the MIP model corresponding to differential characteristics with specific properties
 - > High probability
 - Fixed input/output difference
 - Predefined number of active S-boxes
 - Predefined Hamming weight of the input/output differences

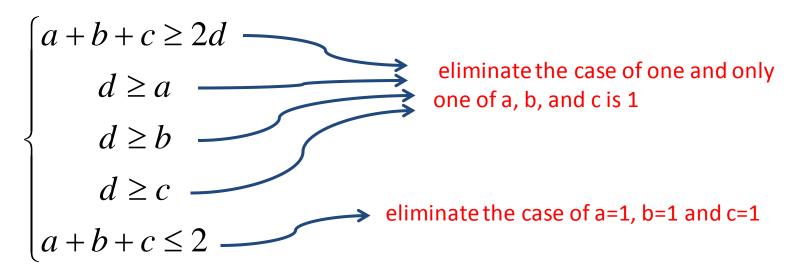
≻...

Constraints imposed on XOR

Constraints imposed on the input and output differences by XOR a

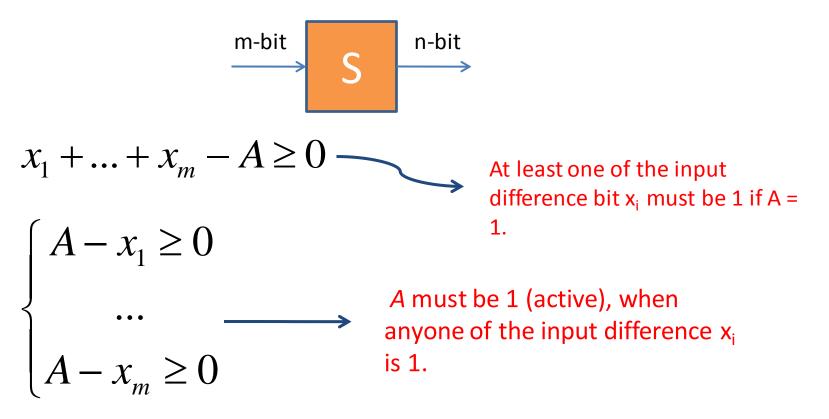


Constraints (where d is a dummy variable and all variables are 0-1)

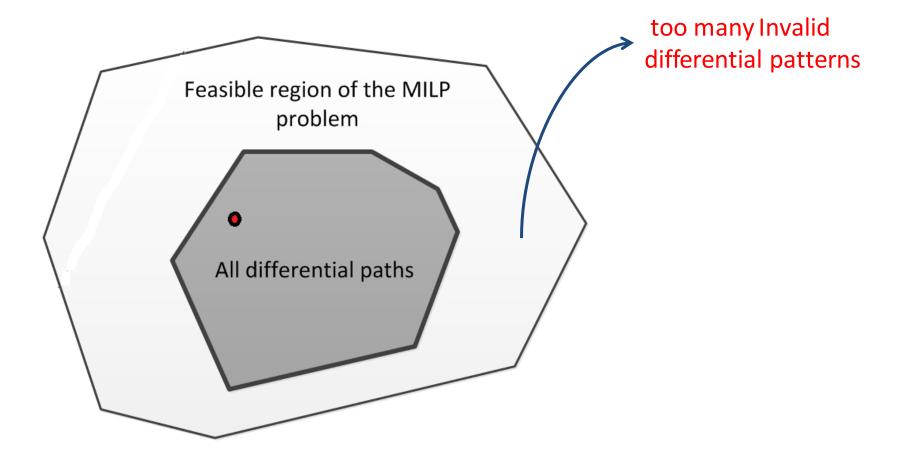


Constraints imposed on the input and output differences by an $m \times n$ S-box (not necessarily invertible)

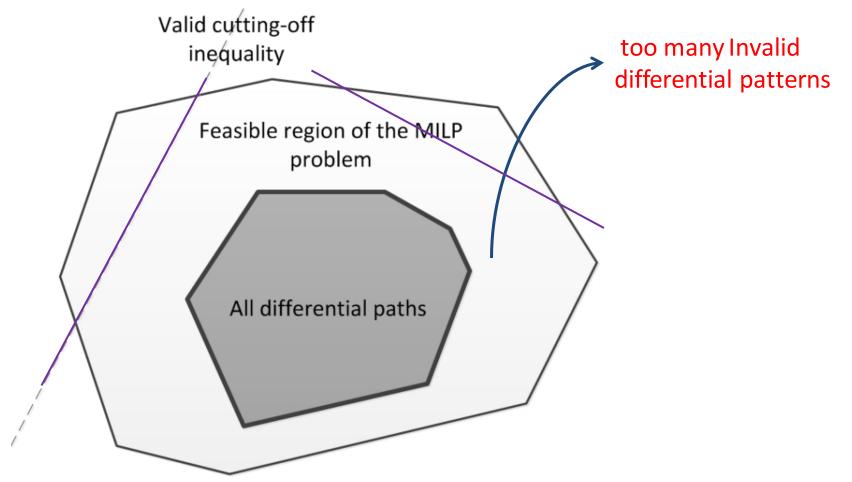
- ✓ Let $x_1, x_2, ..., x_m$ be the input difference, and $y_1, y_2, ..., y_n$ be the output difference
- \checkmark Let A be the variable indicating the activity of the S-box



However, this is too coarse to describe an specific S-box, and result in an feasible region contain many invalid differential patterns



Hence, we need the so called valid cutting-off inequalities to remove some impossible differential patterns of an specific S-box.

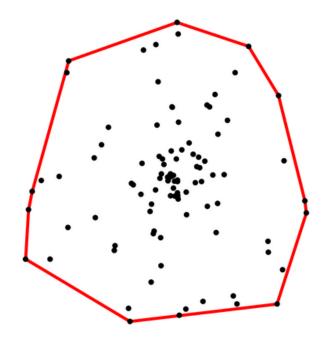


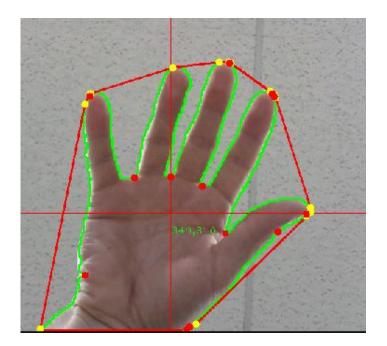
Method for generating valid cutting-off inequalities for an specific S-box

Convex hull computation

Convex hull computation

✓ Convex hull of a set of points in \mathbb{R}^n : the smallest convex set that contains these points.





Convex hull computation

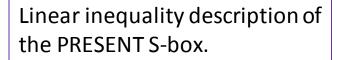
 \checkmark A convex hull can be represented by a set of linear inequalities

Treat the set of all possible differential patterns of an Sbox as a set of points in Rⁿ. For example, the PRESENT S-box:

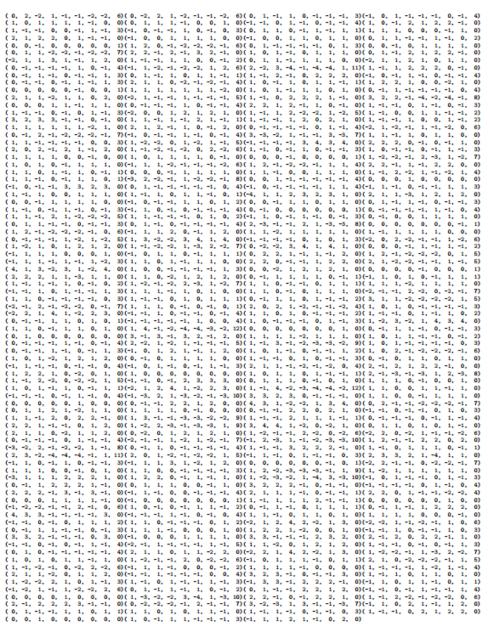
 $\{ (0, 0, 0, 0, 0, 0, 0), (0, 0, 0, 1, 0, 0, 1, 1), (0, 0, 0, 1, 0, 1, 1, 1), \\ (0, 0, 0, 1, 1, 0, 0, 1), (0, 0, 0, 1, 1, 1, 0, 1), (0, 0, 1, 0, 0, 0, 1, 1), \\ (0, 0, 1, 0, 0, 1, 0, 1), (0, 0, 1, 0, 0, 1, 1, 0), (0, 0, 1, 0, 1, 0, 1, 0), \\ (0, 0, 1, 0, 1, 1, 0, 0), (0, 0, 1, 0, 1, 1, 0, 1), \dots \}$

Corresponds to the differential: $0010 \rightarrow 1101$

Then we can compute the linear inequalities representation (H-representation) of the set of differential patterns



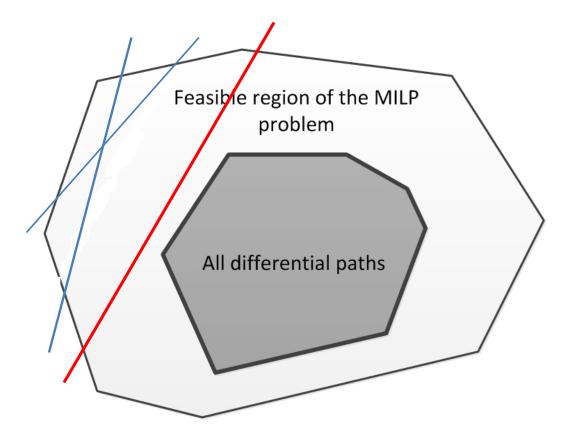
Too many inequalities, which will make the MILP problem too difficult to be solved in practical time



Our method: mixed-integer programming based

Convex hull computation

 \checkmark Can we use less inequalities ? Yes!



Our method: mixed-integer programming based

Convex hull computation

\checkmark Can we use less inequalities ? Yes!

Algorithm 1: Selecting n inequalities from the convex hull \mathcal{H} of an S-box

Input:

 \mathcal{H} : the set of all inequalities in the H-representation of the convex hull of an S-box;

 \mathcal{X} : the set of all possible differential patterns of an S-box;

n: a positive integer.

Output: \mathcal{O} : a set of *n* inequalities selected from \mathcal{H}

- 1 $l^* := None;$
- $2 \ \mathcal{X}^* := \mathcal{X};$

$$3 \mathcal{H}^* := \mathcal{H};$$

$$4 \ \mathcal{O} := \emptyset;$$

- 5 for $i \in \{0, ..., n-1\}$ do
- $l^* :=$ The inequality in \mathcal{H}^* which maximizes the number of removed 6 impossible differential patterns from \mathcal{X}^* ;

7
$$\mathcal{X}^* := \mathcal{X}^* - \{\text{removed impossible differential patterns by } l^*\};$$

8
$$\mathcal{H}^* := \mathcal{H}^* - \{l^*\}$$

- $\begin{array}{c|c} \mathbf{8} & \mathcal{H}^* := \mathcal{H}^* \{l^*\};\\ \mathbf{9} & \mathcal{O} := \mathcal{O} \ \cup \ \{l^*\}; \end{array}$
- 10 end

11 return \mathcal{O}

Applications, Improvements and Adaptations for Specific Ciphers

Applications

- Obtain security bounds w.r.t. (related-key) differential attack
- Search for or enumerate characteristics with specific properties
- Improvements and Adaptations
 - ✓ Search for the best characteristic of some ciphers with small S-boxes, e.g. 4 × 4 S-boxes
 - Construct MIP model whose feasible region is exactly the set of all valid differential characteristics for SIMON like ciphers

 \checkmark

Applications

Obtain security bounds

✓ Set the objective function to minimize the number of active S-boxes. The optimized solution will give a lower bound of the #Active S-boxes

Search for or enumerate characteristics

- ✓ Every feasible solution corresponds to a characteristic
- ✓ After find a solution, add a constraint to the MIP model such that this solution is not feasible any more and find another solution. Repeating this again and again we can enumerate the feasible region

Improvements and Adaptations for Specific Ciphers

Search for the best characteristic of some ciphers with small S-boxes, e.g. 4×4 S-boxes

✓ The idea is to encode the probability information of into the differential pattern of the S-box

Search for the best characteristic of ciphers with 4×4 S-boxes

□ For every differential pattern $(x_0, \dots, x_3, y_0, \dots, y_3)$ of a 4 × 4 S-box, a corresponding differential pattern with probability information can be constructed in the following way

 $(x_0, x_1, x_2, x_3, y_0, y_1, y_2, y_3; p_0, p_1) \in \{0, 1\}^{8+2}$

- $\begin{cases} (p_0, p_1) = (0, 0), \text{ if } \Pr_S[(x_0, \dots, x_{\omega-1}) \to (y_0, \dots, y_{\nu-1})] = 1; \\ (p_0, p_1) = (0, 1), \text{ if } \Pr_S[(x_0, \dots, x_{\omega-1}) \to (y_0, \dots, y_{\nu-1})] = 2^{-2}; \\ (p_0, p_1) = (1, 1), \text{ if } \Pr_S[(x_0, \dots, x_{\omega-1}) \to (y_0, \dots, y_{\nu-1})] = 2^{-3}. \end{cases}$
- □ That is, use the extra two bit (p_0, p_1) to encode the probability information, and the probability of the differential pattern

$$(x_0,\cdots,x_3,y_0,\cdots,y_3)$$

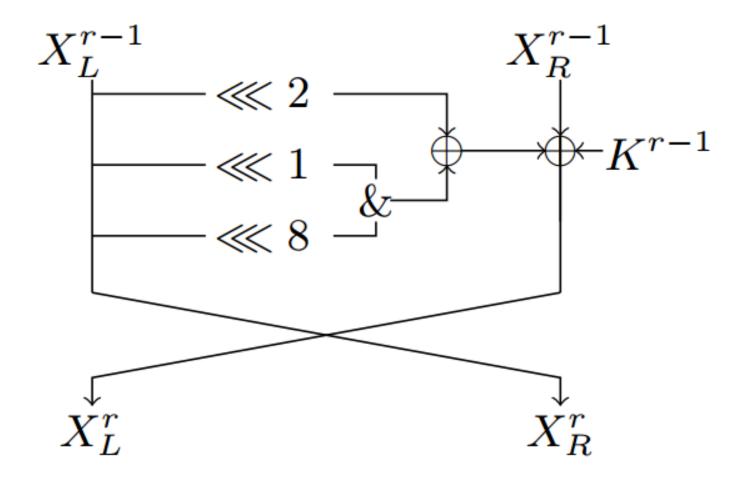
is $2^{-(p_0+2p_1)}$

 \square Set the objective function to be Minimize the sum of all $p_0 + 2p_1$

Improvements and Adaptations for Specific Ciphers

Improvements and Adaptations

- Construct MIP model whose feasible region is exactly the set of all valid differential characteristics for SIMON like ciphers
- ✓ ...



□ The nonlinear layer of SIMON32 can be described by a nonlinear function $F : \mathbb{F}_2^{16} \to \mathbb{F}_2^{16}$ $F(x) = (x <<<1) \cdot (x <<<8), x = (x_0, \cdots, x_{15}) \in \mathbb{F}_2^{16}$

Let $\Delta = (\Delta_0, \dots, \Delta_{15}) \in \mathbb{F}_2^{16}$, and $\delta = (\delta_0, \dots, \delta_{15}) \in \mathbb{F}_2^{16}$, then the differential $\Delta \to \delta$ is a valid if and only if the following system of equation has a solution

$$\begin{cases} \delta_{0} = \Delta_{1} \cdot x_{8} + \Delta_{8} \cdot x_{1} \\ \delta_{1} = \Delta_{2} \cdot x_{9} + \Delta_{9} \cdot x_{2} \\ \delta_{2} = \Delta_{3} \cdot x_{10} + \Delta_{10} \cdot x_{3} \\ \delta_{3} = \Delta_{4} \cdot x_{11} + \Delta_{11} \cdot x_{4} \\ \delta_{4} = \Delta_{5} \cdot x_{12} + \Delta_{12} \cdot x_{5} \\ \delta_{5} = \Delta_{6} \cdot x_{13} + \Delta_{13} \cdot x_{6} \\ \delta_{6} = \Delta_{7} \cdot x_{14} + \Delta_{14} \cdot x_{7} \\ \delta_{7} = \Delta_{8} \cdot x_{15} + \Delta_{15} \cdot x_{8} \end{cases} \begin{cases} \delta_{8} = \Delta_{9} \cdot x_{0} + \Delta_{0} \cdot x_{9} \\ \delta_{9} = \Delta_{10} \cdot x_{1} + \Delta_{1} \cdot x_{10} \\ \delta_{9} = \Delta_{10} \cdot x_{1} + \Delta_{1} \cdot x_{10} \\ \delta_{10} = \Delta_{11} \cdot x_{2} + \Delta_{2} \cdot x_{11} \\ \delta_{11} = \Delta_{12} \cdot x_{3} + \Delta_{2} \cdot x_{11} \\ \delta_{11} = \Delta_{12} \cdot x_{3} + \Delta_{3} \cdot x_{12} \\ \delta_{12} = \Delta_{13} \cdot x_{4} + \Delta_{4} \cdot x_{13} \\ \delta_{13} = \Delta_{14} \cdot x_{5} + \Delta_{5} \cdot x_{14} \\ \delta_{14} = \Delta_{15} \cdot x_{6} + \Delta_{6} \cdot x_{15} \\ \delta_{15} = \Delta_{0} \cdot x_{7} + \Delta_{7} \cdot x_{0} \end{cases}$$

- □ Let $Sol(\delta_0 = \Delta_1 \cdot x_8 + \Delta_8 \cdot x_1)$ be the set of all 0-1 solutions for this equation.
- The vectors $(\delta_0, \Delta_1, x_8, \Delta_8, x_1)$ in $Sol(\delta_0 = \Delta_1 \cdot x_8 + \Delta_8 \cdot x_1)$ are given below

(0,	0,	0,	0,	0)	(0,	0,	0,	0,	1)
(0,	0,	0,	1,	0)	(0,	0,	1,	0,	0)
(0,	0,	1,	0,	1)	(0,	0,	1,	1,	0)
(0,	1,	0,	0,	0)	(0,	1,	0,	0,	1)
(0,	1,	0,	1,	0)	(0,	1,	1,	1,	1)
(1,	0,	0,	1,	1)	(1,	0,	1,	1,	1)
(1,	1,	0,	1,	1)	(1,	1,	1,	0,	0)
(1,	1,	1,	0,	1)	(1,	1,	1,	1,	0)

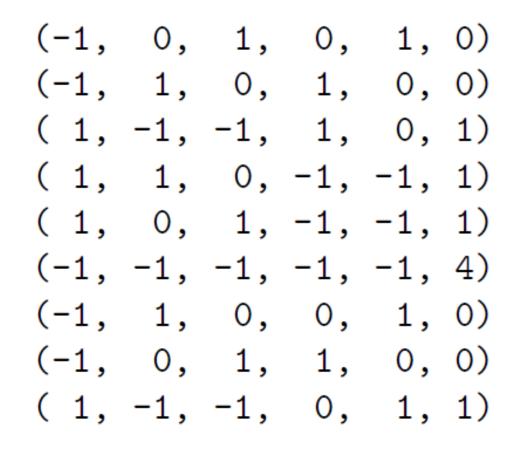
D The H-representation of the convex hull of $Sol(\delta_0 = \Delta_1 \cdot x_8 + \Delta_8 \cdot x_1)$ is given below

$$(0, -1, 0, 0, 0, 1)$$
 $(0, 0, -1, 0, 0, 1)$
 $(0, 0, 0, -1, 0, 1)$ $(-1, 1, 0, 0, 1, 0)$
 $(-1, 0, 1, 0, 1, 0)$ $(-1, 0, 0, 0, 0, 1)$
 $(0, 0, 0, 0, 1, 0)$ $(1, -1, -1, 0, 1, 1)$
 $(0, 1, 0, 0, 0, 0)$ $(0, 0, 0, 0, -1, 1)$
 $(-1, 1, 0, 1, 0, 0)$ $(1, 0, 0, 0, 0, 0, 0)$
 $(-1, 0, 1, 1, 0, 0)$ $(0, 0, 0, 1, 0, 0)$
 $(1, 0, 1, -1, -1, 1)$ $(0, 0, 1, 0, 0, 0)$
 $(1, -1, -1, 1, 0, 1)$ $(1, 1, 0, -1, -1, 1)$
 $(-1, -1, -1, -1, -1, 4)$

where a 6-dimensional vector $(\lambda_0, \dots, \lambda_4, \gamma)$ denotes the linear inequality

$$\lambda_0\delta_0 + \lambda_1\Delta_1 + \lambda_2x_8 + \lambda_3\Delta_8 + \lambda_4x_1 + \gamma \ge 0$$

From the H-representation we can derive the critical set (using the greedy algorithm)



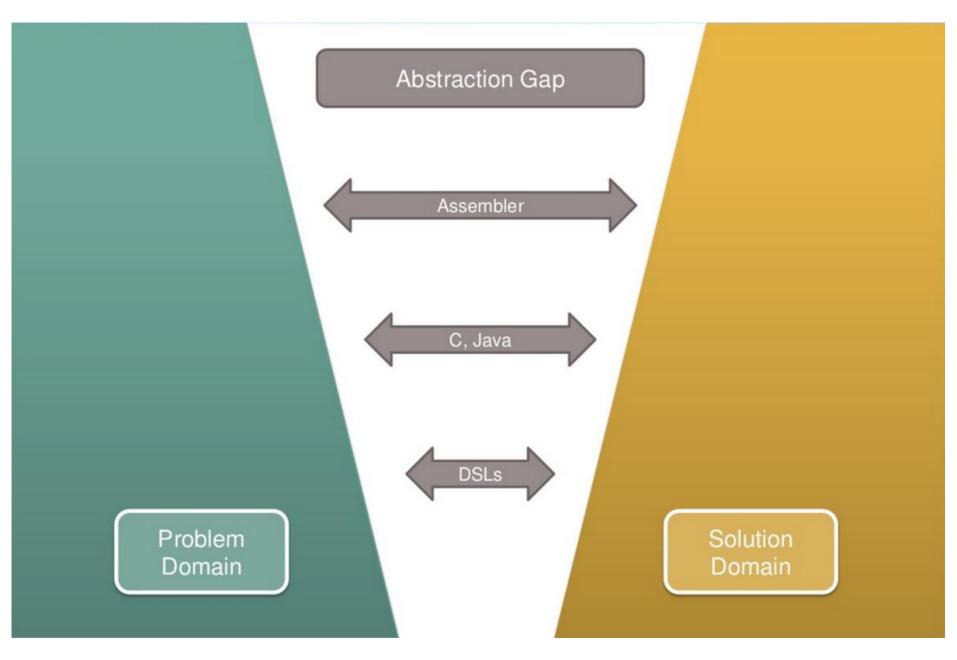
Domain Specific Language (DSL) for Constraint Programming based Automatic Cryptanalysis

Domain Specific Language

A Domain-Specific Language (DSL) is a computer language specialized to a particular application domain.

Compared with General Purpose Language (GPL), DSL offers ...

- ✓ Higher abstractions
- ✓ Avoid redundancy
- $\checkmark\,$ Separation of concerns
- \checkmark Use domain concepts



• http://modeling-languages.com/introduction-to-domain-specific-languages-slides/

Domain Specific Language

Examples

SQL

```
CREATE TABLE Employee (
   id INT NOT NULL IDENTITY (1,1) PRIMARY KEY,
   name VARCHAR(50),
   surname VARCHAR(50),
   address VARCHAR(255),
   city VARCHAR(60),
   telephone VARCHAR(15),
```

```
    HTML
```

```
<html>
<head>
<title>Example</title>
</head>
<body>
Example
</body>
</html>
```

CSS

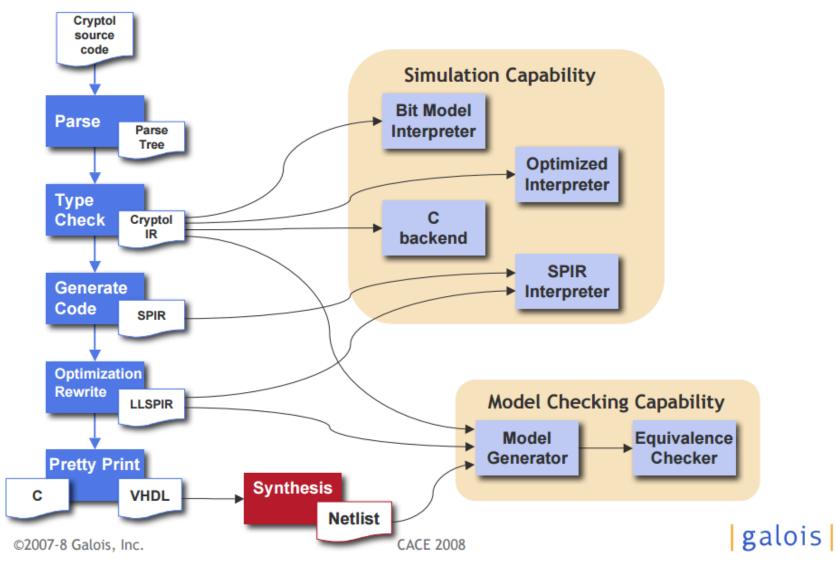
```
body {
   text-align: left;
   font-family: helvetica, sans-serif;
}
h1 {
   border: 1px solid #b4b9bf;
   font-size: 35px;}
```

LaTeX

```
\ifthenelse{\boolean{showcomments}}
 {\newcommand{\nb}[2]{
   \fcolorbox{gray}{yellow}{
    \bfseries\sffamily\scriptsize#1
   }
   {\sf\small\textit{#2}}
   }
   \newcommand{\version}{\scriptsize$-$working$-$}
  }
  {\newcommand{\nb}[2]{}
   \newcommand{\version}{}
}
```

A Domain Specific Language for Crypto Implementation

D Cryptol



Constraint Programming

A mathematical problem involving a set of variables, constraints, and some objective.

□ Constraint programming examples:

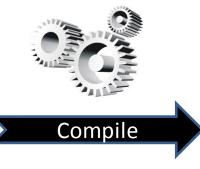
- ✓ SAT problem
- ✓ Mixed Integer Programming
- ✓ Satisfiability Modulus Theory (SMT) problems
- ✓ Constraint Integer Programming

✓ ...

□ Already used in cryptographic research

AES, DES, PRESENT, LBlock, PRIDE, Rectangle ...

Describe the algorithms using DSL (even graphical languages)



MIP, SAT, SMT models for automatic cryptanalysis

Main References

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- 2. Sareh Emami, San Ling, Ivica Nikolic, Josef Pieprzyk and Huaxiong Wang. The Resistance of PRESENT-80 Against Related-Key Differential Attacks. Cryptology ePrint Archive, Report 2013/522, 2013.
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- 11. Finding the Best Characteristics of Some Bit-oriented Block Ciphers and Automatic Enumeration of (Related-key) Differential and Linear Characteristics with Predefined Properties. IACR Cryptology ePrint Archive 2014: 747 (2014)

Thanks!