



*Improving
Counter-cryptanalysis*

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Part I – Weak signature schemes

Part II – Counter-cryptanalysis

Part III – Flame

Part IV – Improvements

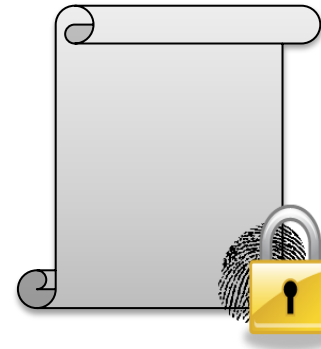
To Conclude...

Introduction Digital Signatures



Digital signature schemes

- One of the pillars for P.K.I.s
- Used to ensure authenticity in/of
 - Browsers
 - Documents
 - Email
 - Software updates
 - Downloadable content
 - Currency transactions
- Hash-Then-Sign:
 {MD5,SHA-1,SHA-2}-{RSA,DSA}
- Hash collision $MD5(A)=MD5(B) \Rightarrow$ forgery

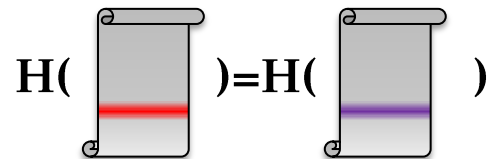




Collision attacks on MD5 & SHA-1

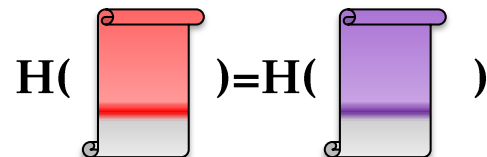
- Distinguish between 2 types
 - Identical prefix

$$H(P | C | S) = H(P | C' | S)$$



- Chosen-prefix

$$H(P | C | S) = H(P' | C' | S)$$



- P, P', S : Free to choose s/t $|P| = |P'|$
- C, C' : Generated based on P and P' , $|C| = |C'| \in [64B, 1KB]$

Introduction

Collision attacks



	MD5		SHA-1		SHA-256	
	Id.Pr.	Ch.Pr.	Id.Pr.	Ch.Pr.	Id.Pr.	Ch.Pr.
Birthday	$2^{64.3}$	$2^{64.8}$	$2^{80.3}$	$2^{80.8}$	$2^{128.3}$	$2^{128.8}$
2004	2^{40}		2^{69}			
2005	2^{37}		(2^{63})			
2006	2^{32}	2^{49}				
2007	2^{25}	2^{42}	(2^{61})			
2008	2^{21}					
2009	2^{16}	2^{39}				
2010						
2011						
2012			2^{61}	2^{77}		
today	2^{16}	2^{39}	2^{61}	2^{77}	$2^{128.3}$	$2^{128.8}$

Published collision attacks on MD5 & SHA-1

Introduction

Collision attacks



Notes

- Generate your own MD5 chosen-prefix collision attack in a day using Project HashClash:
<https://code.google.com/p/hashclash/>
- No publicly known collision for SHA-1 has been found yet
- First SHA-1 collision more likely to be constructed by nation-states than academia due to required resources, see:
http://www.schneier.com/blog/archives/2012/10/when_will_we_see.html

Introduction Collision strategies



Strategies for *meaningful* colliding files

- Using identical-prefix collisions

- Meaningful C and C'

Hard

C = "... of money is \$10,000.00..."

C' = "... of money is \$20,000.00 ..."

C = "... OFFSET=X ..."

C' = "... OFFSET=Y ..."

- IF-THEN-ELSE construct

IF (C ==C) THEN ... ELSE ...

IF (C'==C) THEN ... ELSE ...

Easy, but requires
IF-THEN-ELSE

- Using chosen-prefix collisions

- Meaningful different P, P' & hide C and C' in message

Easy

P = "I owe you \$20"

C = <hidden image>

P' = "You will inherit all my possessions"

C' = <hidden image>

Introduction

Abuse examples



- Identical-prefix
 - Colliding Software [Kam04,Mik04]
 - Colliding PostScript documents [DL05,GIS05]
 - Colliding X.509 certificates (same ID, diff. RSA moduli) [LdW05]

- Chosen-prefix
 - Colliding PDF documents [SLdW07]
 - Colliding Software [SLdW07]
 - Colliding X.509 certificates (diff. IDs) [SLdW07]
 - Rogue Certification Authority [SSALMOdW09]
 - **Rogue Windows Update signing certificate [Flame12]**

Introduction

Solutions for weak signatures schemes



What to do when a signature scheme is broken?

The easy answer: “migrate to a more secure scheme”
i.e., move from MD5-RSA to SHA-2-RSA

Who should migrate?

Signers: generate SHA-2-RSA signatures

Problems: compatibility / deployment issues, risk-cost trade-off, human,...

Result: forgeries can still be constructed till the last signer migrates

Verifiers: don't accept MD5-RSA signatures

Problem: too many old signatures in use to just invalidate them all at once

Result: old and new forgeries can be abused against nearly everyone

The easy answer is not a practical solution for the near future



Our answer: detect forged signatures

Verifiers: don't accept *forged* MD5-RSA signatures

Results:

- Old legitimate signatures are still valid
- Verifier protected against forgeries
- Independent of migration by signers

How to detect forged signatures?: counter-cryptanalysis!



Part II – Counter-cryptanalysis

Part III – Flame

Part IV – Improvements

To Conclude...

Counter-cryptanalysis New paradigm



New paradigm: counter-cryptanalysis

- Strengthen weak cryptographic primitives
- Detect cryptanalytic attacks at the cryptographic level
- Counter-cryptanalysis in principle enables the continued secure use of weak cryptographic primitives
- No strengthened redesign \Rightarrow no compatibility issues
- May be used during migration to strengthened redesigns in the real world

Why is that possible?

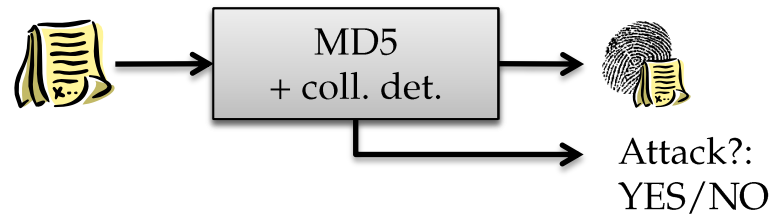
- Dedicated cryptanalytic attacks are highly specialized
- Active attacks may introduce subtle unavoidable anomalies
- Similar cryptanalytic techniques can be used to detect those anomalies
- This approach may detect an entire class of attacks that all introduce the same unavoidable anomalies

Counter-cryptanalysis Collision detection



First practical example: collision detection

- Detect whether message was constructed using collision attack
- Single message of collision pair sufficient
- Application to MD5 & SHA-1
- Computational cost
 - MD5 factor x 224
 - SHA-1 factor x 15
 - Much less using early-abort: WIP
- Based on crucial properties of the known cryptanalysis on MD5 & SHA-1
 - Attacks exploit trivial differential steps with probability (close to) 1 to be able to obtain 'low' complexity
 - Very few message block differences result in attacks with 'low' complexity

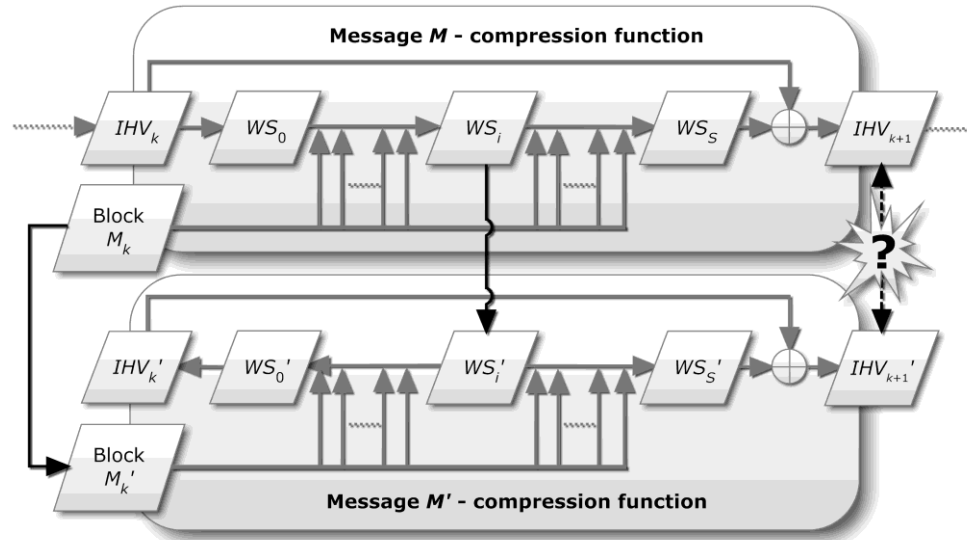


Counter-cryptanalysis Collision detection



Basic algorithm: detect last near-collision block

1. Guess message block difference & difference at trivial step i
2. Determine M_k' from M_k and WS_i' from WS_i
3. Reconstruct computation
4. Check whether collision in chaining value is obtained



If guess was correct then collision is detected with certainty

If guess was incorrect then a false positive occurs with probability $\approx 2^{-N}$

Counter-cryptanalysis Reference implementation



Reference implementation to detect collision attacks

- Available at <http://marc-stevens.nl/research> (at the bottom)
- Library interface to replace existing MD5/SHA-1 implementation
 - *MD5Init/MD5Init_unsafe, MD5Update, MD5Final*
 - *SHA1Init/SHA1Init_unsafe, SHA1Update, SHA1Final*
 - *{MD5,SHA1}Final* returns non-zero value if an attack is detected
 - *{MD5,SHA1}Init_unsafe* always results in correct (and possibly unsafe) hash
 - *{MD5,SHA1}Init* results in correct hash **if no attack has been detected**, otherwise a safe hash is returned
- Command line program
 - `detectcollv <files>`



Anomaly detection for digital signatures

- Online: active protection
 - Signer: protection against malicious signature requests
 - Verifier: protection against forged signatures
 - E.g., for TLS/SSL, OSs (drivers, executables, updates), etc.
- Offline: forensic analysis
 - Main example: spyware Flame



Part III – Flame

Part IV – Improvements
To Conclude...

Flame Overview of Flame



Cf. [Kas12,Sot12]

- Highly advanced malware
- Targeting the Middle-East
- Discovered in May 2012
- Active since 2007 or earlier
- Uncharacteristic features for malware
 - Up to 20 modules: each carefully selected prior to infection
 - Almost 20MB: includes Lua VM & libraries for compression, database, ...
 - Did not spread wildly & evaded discovery for ~5 years
 - Surgical-precision attacks: each target carefully selected
 - Spread itself illegitimately using the Windows Update platform
 - First cryptanalytic attack on hash function found in the 'wild'
 - Developed new variant cryptanalytic attack to do so...



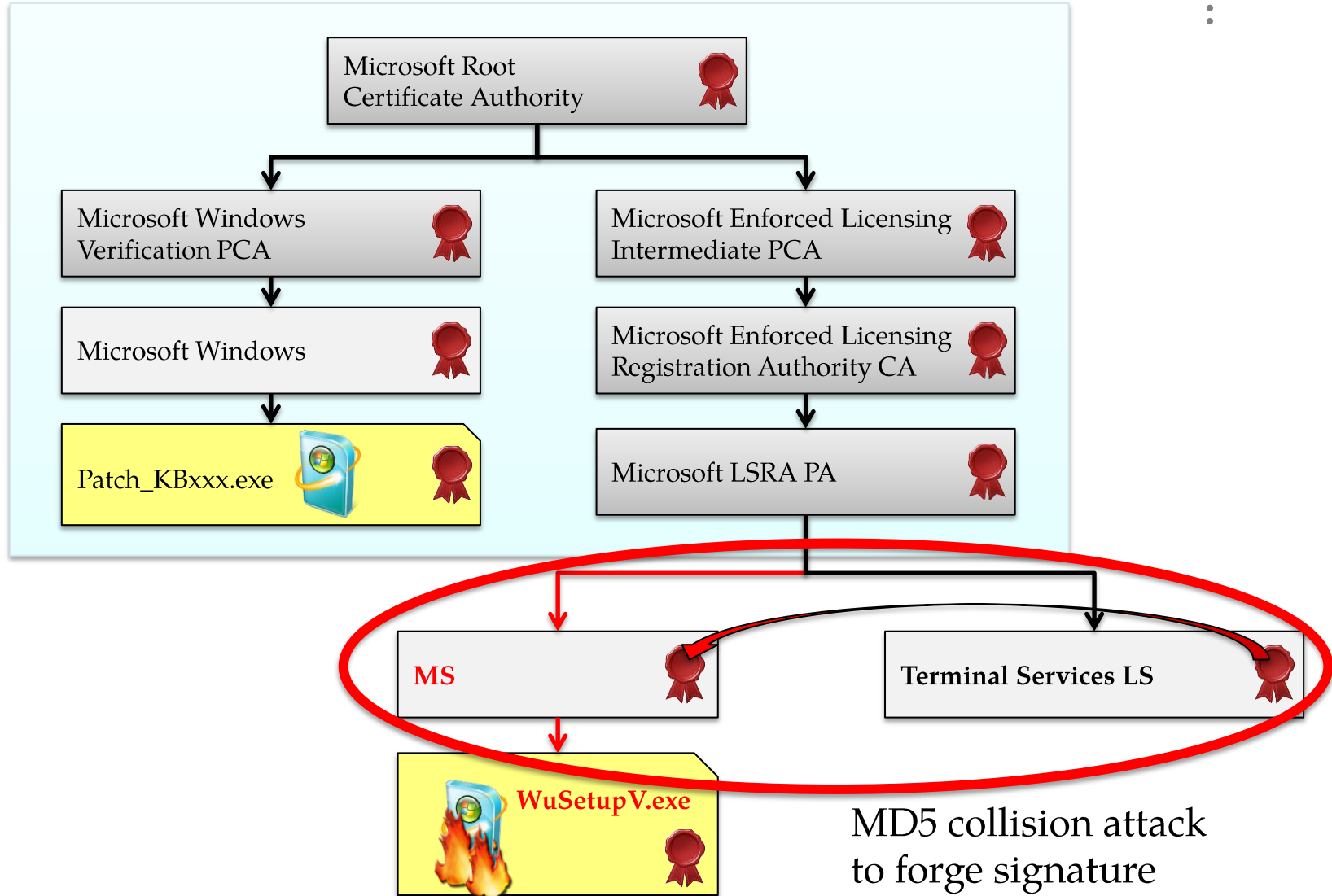
Source: Kaspersky Lab

Flame Propagation



- Man-in-the-middle attack on Windows Update
- Local network attack
 - Registers itself as proxy server for **update.microsoft.com** using WPAD (Web Proxy Auto-Discovery)
 - Windows Update falls back to insecure HTTP
 - depends on digital signatures for security
 - no need to subvert TLS/SSL connection
- Propagation
 - Flame serves fake 'security update' using Windows Update platform
 - Requires properly-signed 'security update'
 - **Uses illegitimate sub-CA** valid since **Feb 2010**
 - ⇒ sub-CA invalid before that time
 - ⇒ this attack was almost certainly done around Feb 2010 or later

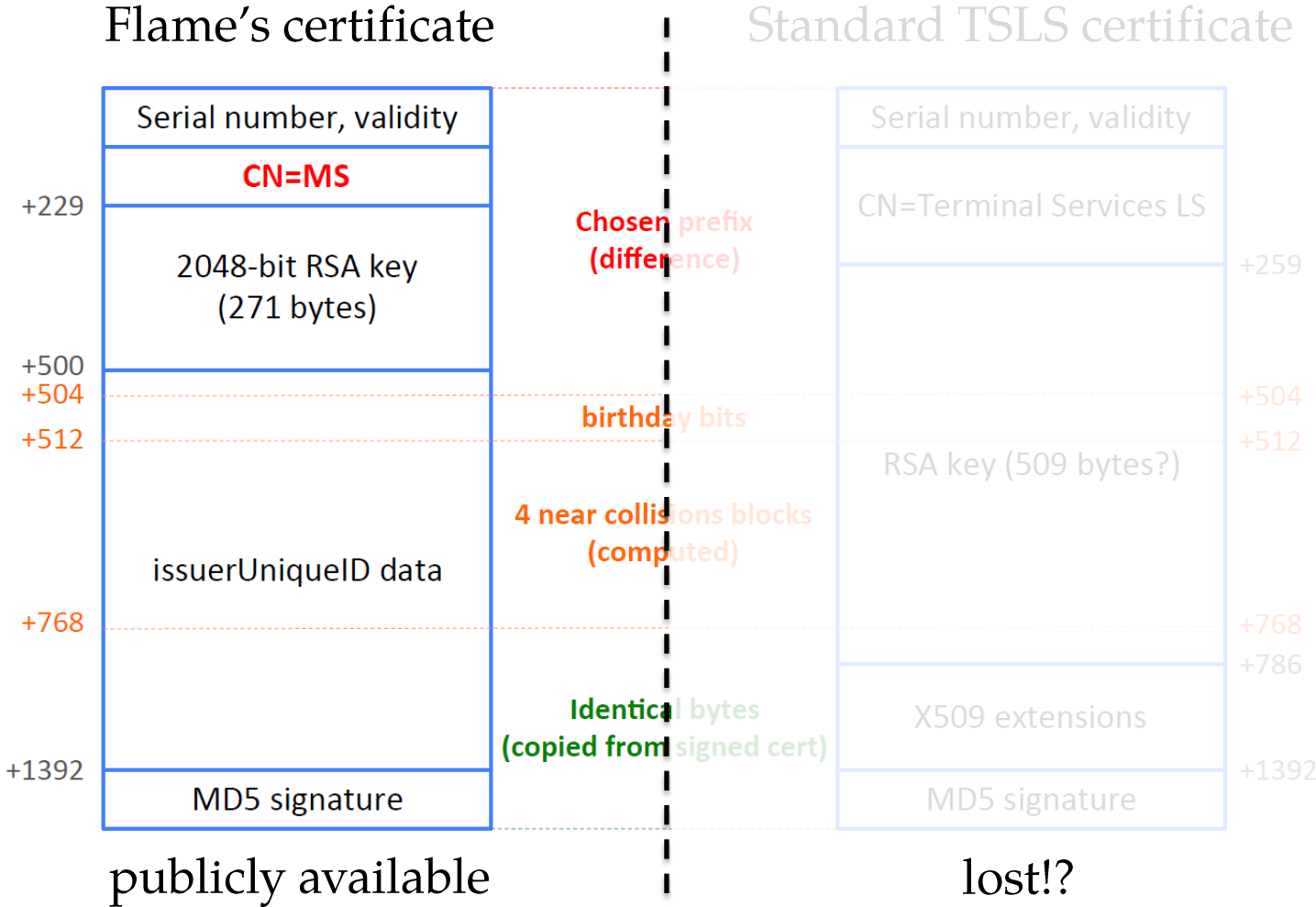
Flame Certificate hierarchy



Flame Chosen-prefix collision attack



Uses chosen-prefix collision attack [SLdW07]:



Source: Alex Sotirov

Flame Analyzing Flame's attack



- Only Flame's "MS" sub-CA certificate public
- The colliding "TSL" certificate is not public (lost?)
- First example for counter-cryptanalysis
 - Assumed chosen-prefix collision attack
 - Only 1 of the 2 colliding certificates available to us
 - Ran proof-of-concept implementation (from 2008)
 - chosen-prefix collision detected
 - 4 near-collision blocks recovered
 - all differential paths reconstructed
 - <0.03 seconds
 - Differential paths expose use of new variant attack

Flame Recovered differential paths



```

dm4=[[!31!]] dm11=[[!15!]] dm14=[[!31!]]
Q-3: |..... |..... |..... |..... |..... |
Q-2: |00..... |1.1.01. |...1.+ |...-10.. |
Q-1: |110+...1 |.1-.00. |.+... |...-110.. |
Q0: |+-100..0 |.-0+^++1 |.0.+0.11 |.110+... | ok p=1
Q1: |0+--+...- |.-0+++0 |011-0..1 |110+++... | ok p=0.49707
Q2: |+0-0-..00 |.-++00+- |0-1-+.1+ |1+-0+++^ | ok p=0.166016
Q3: |+010-000 |.-+++0+1 |+-+.^+1+ |+-+---+ | ok p=1
Q4: |-00-10+.. |.11-+-0+ |+++11--0 |101-+0. | ok p=1
Q5: |0-+-+--^ |^0110+1- |110+0-0 |0001+1^ | ok p=1
Q6: |+-+---+ |---+--- |-----+ |+++++++ | ok p=1
Q7: |111.-111 |1101011. |110-1001 |+0100.00 | ok p=1
Q8: |00+0.111 |10111101 |1101100. |1110011 | ok p=0.170898
Q9: |..0.1... |.....- |0.10+... |0-...0. | ok p=0.563477
Q10: |..0^...1 |^...0... |0^0-1... |.1....+ | ok p=0.121094
Q11: |..0-...1 |+...-... |.+01... |0..^1. | ok p=0.899414
Q12: |.1-1..^+ |1...+... |.+0... |...+1. | ok p=0.946289
Q13: |.0+1...+ |1...0... |100...1 |...0... | ok p=0.655273
Q14: |.-+...1. |...1... |1.+...1 |...1... | ok p=0.578125
Q15: |.0+...10 |.....- |0.....- |.....- | ok p=0.989258
Q16: |.1+... |0..... |^..... |..... | ok p=0.887695
Q17: |.1..... |.1..... |0..... |^..... | ok p=1
Q18: |.0..... |+...1. |..... |..... | ok p=0.998047
Q19: |..... |..... |..... |..... | ok p=0.864258
Q20: |0..... |^..... |..... |..... | ok p=1
Q21: |0..... |..... |^..... |..... | ok p=0.501953
Q22: |..... |..... |..... |..... | ok p=0.517578
Q23: |..... |..... |..... |..... | ok p=1
Q24: |^..... |..... |..... |..... | ok p=1
-----
Q25-32: |..... |..... |..... |..... | ok p=1
-----
Q33: |0..... |..... |..... |..... | ok p=1
Q34: |1..... |..... |..... |..... | ok p=0.507812
-----
Q35-59: |±..... |..... |..... |..... | ok p=1
-----
Q60: |+.11110. |..... |..... |..... | ok p=1
Q61: |+.11000. |..... |.001.00. |..... | ok p=1
Q62: |.-+--- |..... |0..... |..... | ok p=0.426758
Q63: |.+?0??+. |..... |.-+.+. |..... | ok p=0.855469
Q64: |+.....+ |+++++ |.-.-+. |.....+- |

```

differential path 1

```

dm4=[[!31!]] dm11=[[!-15!]] dm14=[[!31!]]
Q-3: |+..... |..... |..... |..... |..... |
Q-2: |-1...+. |.1.1.0.. |0....1+ |.-+...0. |
Q-1: |+01-..+1 |.0-+.0^ |011+--1 |--+0.10 |
Q0: |1-0-1.+0 |^-0+1+-1 |1011+-0 |001.1^-1 | ok p=0.749023
Q1: |10-.01.+ |+++0+10 |--+111+- |+--0+1- | ok p=0.425781
Q2: |.01.-011 |00+--+0+ |0-+.-0 |++10+0+0 | ok p=0.492188
Q3: |.1.-+11 |+001+^+ |01-+0110 |0+1+0+0+ | ok p=0.833008
Q4: |..-1-11 |+1-+++ |1111-+- |++0+--1 | ok p=1
Q5: |^1^+1-- |10-01011 |0+10-1-+ |0-+++000 | ok p=0.499023
Q6: |+++++ |+++--- |-----+ |--+--- | ok p=1
Q7: |0010-000 |01111011 |1011-111 |10.10010 | ok p=1
Q8: |00000100 |1111111+ |1001111 |1-010111 | ok p=0.672852
Q9: |...-1... |.....- |0..1+... |.1....^ | ok p=0.495117
Q10: |...0...0 |^0...0 |1..+0... |0.....- | ok p=0.895508
Q11: |..0...^ |1...^ |...01... |.....-1 | ok p=0.807617
Q12: |.001...+ |0..... |01..... |.....1. | ok p=1
Q13: |.1-1.0- |1...0. |1^1...1 |...1... | ok p=1
Q14: |.-+...10 |.....0. |1-+...1 |...1... | ok p=0.586914
Q15: |.0+...0 |..... |+01...+ |..... | ok p=0.994141
Q16: |.^+... |0..... |^..... |..... | ok p=0.879883
Q17: |.1..... |.1..... |0..... |^..... | ok p=1
Q18: |.0..... |.-...1. |..... |..... | ok p=0.999023
Q19: |..... |..... |..... |..... | ok p=0.895508
Q20: |0..... |^..... |..... |..... | ok p=1
Q21: |0..... |..... |^..... |..... | ok p=0.487305
Q22: |..... |..... |..... |..... | ok p=0.508789
Q23: |..... |..... |..... |..... | ok p=1
Q24: |^..... |..... |..... |..... | ok p=1
-----
Q25-32: |..... |..... |..... |..... | ok p=1
-----
Q33: |1..... |..... |..... |..... | ok p=1
Q34: |0..... |..... |..... |..... | ok p=0.507812
-----
Q35-59: |±..... |..... |..... |..... | ok p=1
-----
Q59: |+..... |..... |0..... |..... | ok p=1
Q60: |+.....0. |..... |...1001. |110..... | ok p=0.506836
Q61: |-...100 |...0... |...1.1. |00+... | ok p=0.749023
Q62: |+...1- |..... |.-+++ |+-... | ok p=0.948242
Q63: |+...+- |+...+ |...??- |?+... | ok p=0.261719
Q64: |+...-- |+... |.-...- |.-+...+ |

```

differential path 2

Flame Recovered differential paths



```

dm4=[[!31!]] dm11=[[!15!]] dm14=[[!31!]]
Q-3: |.....|
Q-2: |.1.10100 .....11. 10..... .0.....|
Q-1: |^0.0101- .1.0^10. 11.0.... .1.100^|
Q0: |++1-++++ 1001--- --.1.... .1+.110-| ok p=1
Q1: |0-111110 1-1+1+^ -+1.... .01^++-0| ok p=0.96875
Q2: |10-01110 +++1---+ +10+.... 0-0+++++1| ok p=0.374023
Q3: |-0-01^1+ +0+1--10 0-+++^0 01+0+00.| ok p=1
Q4: |--0++-00 0-0+11++ +-+1-+10 -+00+-1.| ok p=1
Q5: |-1++-0-1 +1-00+1- +0+++110- -1--1+^| ok p=1
Q6: |+-+----+ -+++-+--- -++---+ ++++++| ok p=1
Q7: |1000-010 00.1010. 101-0101 +0001.00| ok p=1
Q8: |11+1.101 01011100 -1000101 .1000011| ok p=0.0566406
Q9: |..0.1... ..-... 0.10+... 0-...0.| ok p=0.573242
Q10: |..0^...1 ^.....0. 0^0-1... .1....+| ok p=0.120117
Q11: |..0-...1 +.....- .+01... .0.^1.| ok p=0.889648
Q12: |.1-1.^+ 1....+... .0+0.... .+1.| ok p=0.948242
Q13: |.0+1..-+ 1....0. 100....1 ....0...| ok p=0.631836
Q14: |.-+...1. ....1... 1.+...1 ....1...| ok p=0.585938
Q15: |.0+...10 .....-0.....- .....-| ok p=0.993164
Q16: |.1+..... .0..... .^.....| ok p=0.868164
Q17: |..1..... .1.....0. ^.....^ .....^| ok p=1
Q18: |..0..... .+...1. ....| ok p=0.999023
Q19: |.....- .....- .....| ok p=0.868164
Q20: |0..... .^.....| ok p=1
Q21: |0..... .^.....| ok p=0.495117
Q22: |-.....- .....- .....| ok p=0.509766
Q23: |.....- .....- .....| ok p=1
Q24: |^.....- .....- .....| ok p=1
-----
Q25-32: |.....- .....- .....| ok p=1
-----
Q33: |1.....- .....- .....| ok p=1
Q34: |1.....- .....- .....| ok p=0.493164
-----
Q35-59: |±.....- .....- .....| ok p=1
-----
Q60: |-.....0. ....1. ....| ok p=1
Q61: |-.0110.0 ..... .1....0. ....| ok p=0.514648
Q62: |+.01+. .... .0....+ .....| ok p=0.492188
Q63: |+.+---?- .....+ .....| ok p=0.395508
Q64: |+.+.- .....+ +.....- .....|

```

differential path 3

```

dm4=[[!31!]] dm11=[[!-15!]] dm14=[[!31!]]
Q-3: |+.....|
Q-2: |+.....0+. .... 000+--- .000..1|
Q-1: |+.....+. 11....++ ++1101+. 10011..1|
Q0: |001.1+- .01^..111 -+-----0 11+-+11-| ok p=1
Q1: |011.0.+ .-+^++1+ ++0000-1 +-+0-11+| ok p=0.742188
Q2: |+---0- .-+1+0--0 1+1-1-++ -1-00+-| ok p=0.756836
Q3: |+-+1-^1. .+100-- 10---1+0 ---0+-+1| ok p=1
Q4: |-010+-1. 10-1-01+ 0-000-1- 0+-10-1-| ok p=1
Q5: |+00-+00^ 0+-+11-0 ++0-0-111 01+-+100| ok p=1
Q6: |+++++ +---+ -+---+ -+---+| ok p=0.506836
Q7: |.111-110 01.010.0 0101-110 1101.011| ok p=0.735352
Q8: |11110110 0101000+ -0101111 0-100111| ok p=0.0507812
Q9: |...-1... .-...1 0..1+... .1....^| ok p=0.522461
Q10: |...0...0 ^0.....0 1..+0... .0.....-| ok p=0.895508
Q11: |..0...^0 -1...^1. ....01.....1| ok p=0.822266
Q12: |.001..-+ 0.....- .111.....1..| ok p=1
Q13: |.1-1.0- 1....0. 100....1 ....1...| ok p=1
Q14: |.-+...10 .....0.. 1-+...1 ....1...| ok p=0.556641
Q15: |.0+...0 .....+01.....+ .....| ok p=0.998047
Q16: |.^+..... .0..... .^.....| ok p=0.892578
Q17: |..1..... .1.....0. ^.....^ .....^| ok p=1
Q18: |..0..... .-...1. ....| ok p=0.999023
Q19: |.....- .....- .....| ok p=0.860352
Q20: |0..... .^.....| ok p=1
Q21: |0..... .^.....| ok p=0.485352
Q22: |+.....- .....- .....| ok p=0.501953
Q23: |.....- .....- .....| ok p=1
Q24: |^.....- .....- .....| ok p=1
-----
Q25-32: |.....- .....- .....| ok p=1
-----
Q33: |0.....- .....- .....| ok p=1
Q34: |1.....- .....- .....| ok p=0.498047
-----
Q35-59: |±.....- .....- .....| ok p=1
-----
Q60: |+.....0. ....00. ....| ok p=1
Q61: |+.....1. .... 11....1. ....| ok p=0.525391
Q62: |-.....- ..... 10...+ .....| ok p=0.493164
Q63: |+.....- .....+ .....?-. .....| ok p=0.50293
Q64: |...-+... .....+ .....- .....+|

```

differential path 4

Flame

Yet unknown collision attack



Yet unknown chosen-prefix collision attack

1. Other differential path family

- Same message differences for all 4 near-collision attacks (up to sign)
 $\delta m_4 = 2^{31}$, $\delta m_{11} = \pm 2^{15}$, $\delta m_{14} = 2^{31}$ (same as [WY04])
- No systematic elimination using $\delta m_{11} = \pm 2^b$ as in [SSA+09]

2. Yet unknown birthday search

- Birthday search is preprocessing phase to find differences that can be cancelled by differential path family
- Less flexible differential path family \Rightarrow Higher complexity birthday search
- Approx. 2^{49} MD5 compression function calls
- To compare: our attack in total has average complexity of $2^{44.6}$ MD5 calls

Flame

Yet unknown collision attack



Yet unknown chosen-prefix collision attack

3. Yet unknown differential path construction algorithm

- Differences in *all* bit positions of Q6 in *all* 4 near-collision attacks
- Not a characteristic of known construction algorithms
- Their connection search-space strictly contained in our search-space
- Initial tests show this approach to be significantly slower than our approach
- Probability successful connection over 4 steps drops approx. by factor 2^{-13}

For a more extensive analysis, see <http://eprint.iacr.org/2013/358>



Part IV – Improvements

To Conclude...

Improving counter-cryptanalysis

False negatives



Improving counter-cryptanalysis

- Reduce chance at false negatives
- Need 'exhaustive' list of $(\delta B, \delta WS)$

- Case SHA-1
 - Need 'exhaustive' list of feasible disturbance vectors
 - Heuristic searches have found only two interesting classes of D.V.s
 - Make selection based on cost function, e.g. [Ste12]

- Case MD5
 - Finding feasible δB unfinished work
 - Literature focused on finding attacks better than Wang et al.'s
 - W.I.P. heuristic 'exhaustive' surveys for δB
 - Using cost function determining lower-bound for attack complexity

Improving counter-cryptanalysis

Early stop



Improving counter-cryptanalysis

- Each $(\delta B, \delta WS)$ -guess costs 1 full compression function call
- Speed-up by early stop
 - Use very fast pre-check and only do full work with low probability
 - Without introducing possible false negatives
 - Find unavoidable conditions on message and state bits
(conditions necessary for all possible feasible attacks based on $(\delta B, \delta WS)$)
- Case SHA-1
 - Determine unavoidable message bitrelations over steps 30-70
 - E.g., using exact joint local-collision analysis [Ste12]
 - Expect at least 4 bitrelations per Disturbance Vector
 - Would result in total cost factor < 2 instead of 16

Improving counter-cryptanalysis

Early stop



Improving counter-cryptanalysis

- Case MD5 (per $\delta B, \delta W S_i$)
 - More complex, multiple types of pre-checks
 - Do check whether transitions $(0,0,0,0) \leftrightarrow (2^{31}, 2^{31}, 2^{31}, 2^{31})$ happen correctly
 - Determine lower-bound on attack complexity $< 2^{64}$
 - Determine $K = \#$ last steps that may vary s/t complexity $< 2^{64}$
 - Do check whether $(2^{31}, 2^{31}, 2^{31}, 2^{31}) \rightarrow (2^{31}, 2^{31}, 2^{31}, 2^{31})$ happen correctly for first $16-K$ steps of round 4
 - Find unavoidable bitconditions in round 2 & 3
 - i.e., check whether forcing negated bitcondition results in complexity $\geq 2^{64}$



To conclude...

Conclusions



- Migrate away from MD5 and SHA-1 based signature schemes
- The easy answer of “migrate to more secure signature schemes” is not a practical solution for the near future
- Instead allow old signatures, but protect verifiers against forgeries
- Real-time signature forgery detection possible
 - works for collision attacks on MD5 & SHA-1 [Ste12]
 - recovers full differential paths
- Reference implementation available
 - Feedback requested!
 - Let me know where it is used!

Conclusions



- Flame uses chosen-prefix collision attack ‘in the wild’
 - But an entirely new variant!
 - Different differential path family than [SSA+09]
 - Yet unknown birthday search
 - Yet unknown block-wise elimination procedure
 - Yet unknown differential path construction algorithm
 - New attack has higher complexity than [SSA+09]

Open questions



- Who made Flame?
 - Evidence points to world-class cryptanalysts, not just hackers
 - Adds to predominant speculation of nation-state behind Flame

- Why develop a new variant attack before Feb 2010?
 - But our attack implementation is public since June 2009 (see [Ste12])
 - Requires large effort, done in parallel
 - Nevertheless: exposes their cryptanalytic knowledge

Open questions



- Will the first successful SHA-1 attack be due to scientific efforts or not?
 - Recent years have shown almost no public efforts on SHA-1.
 - Flame's attack on MD5 was developed independently and in parallel to public scientific efforts.
 - Perhaps attacks on SHA-1 are as well...
 - Nation-states have more computing resources than academics...

- Counter-cryptanalysis
 - New paradigm against cryptanalytic attacks
 - Collision detection first practical example
 - Can we construct other (practical) examples against known cryptanalytic attacks?



Thank you for your attention

Questions?

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